## RAMCloud Design Review

# Indexing

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April 1, 2010

#### Introduction

- Should RAMCloud provide indexing?
  - o Leave indexes to client-side using transactions?
- Many apps have similar indexing needs
  - Or compose standard mechanisms to suit their needs
  - o Can optimize for common needs on server-side

#### Implementation Issues

- Indexing on "opaque" data
- Partitioning Indexes
- Consistency
- Recovery/Availability of Indexes

- Problem: RAMCloud treats objects as opaque
  - Server-side indexing without understanding the data?

```
put(tableId, person.objectId, person.pickle())
```

Blob

- Problem: RAMCloud treats objects as opaque
  - Server-side indexing without understanding the data?
- Idea: Apps provide search keys explicitly
  - Apps understand the data

- Problem: RAMCloud treats objects as opaque
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- Can eliminate redundancy
  - Search keys need not be repeated in object
  - Search keys + Blob are returned to app on get/lookup

Lookups are distinct from gets

```
lookup(tableId, `last', `Power')
```

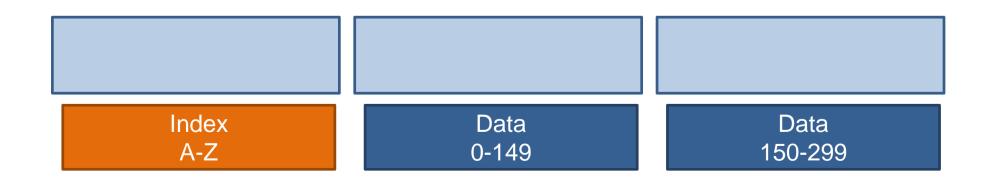
- Put atomically updates indexes and object
  - Details to follow

#### **Partitioning Indexes**

Co-locate index and data



- Large tables?
- Large indexes?
  - Can't avoid multi-machine operations



#### **Partitioning Indexes**

Split indexes on search key



One extra access per lookup and put

#### Split indexes on object ID



- Lookups go to all index fragments
- Puts are always local
- Ordered enumeration of the index is problematic

#### **Partitioning Indexes: Thoughts**

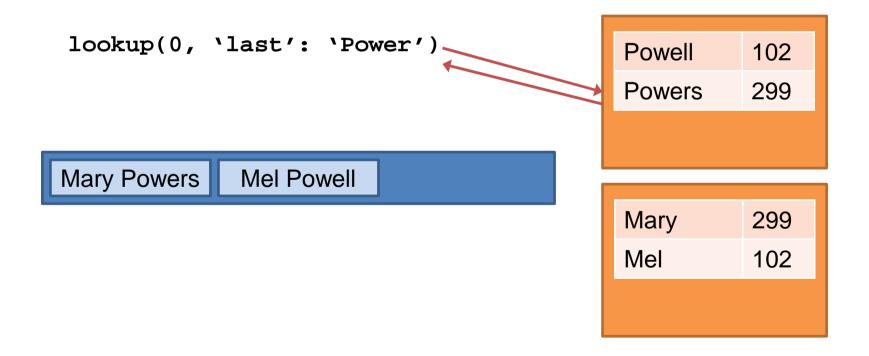
- Our decision (for now): On search key
  - 1. Don't want weakest-link lookup performance
  - 2. To support enumerate and cursors for range queries

#### Consistency

- Problem: Index/Object inconsistency on puts
  - Since object and index may reside on different hosts
  - Apps may see index entries for objects not yet written
- Avoid fancy commit protocol, if possible
- Idea: Index entries "commit" on object put
  - Object puts are atomic
  - Index entries invalid until corresponding put finishes

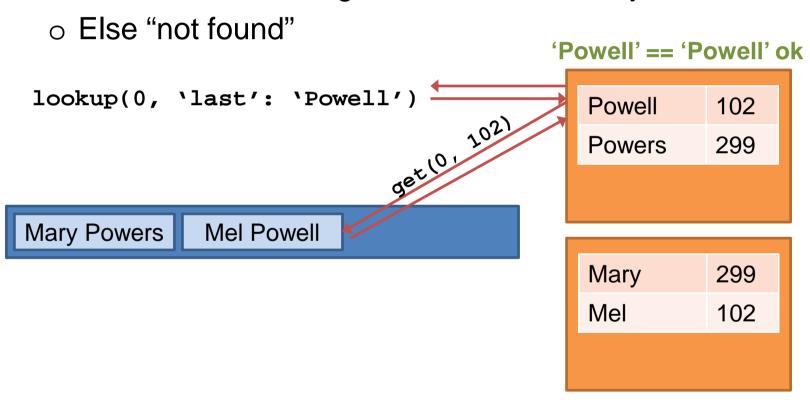
#### **Consistency: Lookup**

- Request goes directly to correct index partition
  - o "Not found" returns immediately



#### **Consistency: Lookup**

- Consistency is checked on hit
  - If table and index agree the return the object



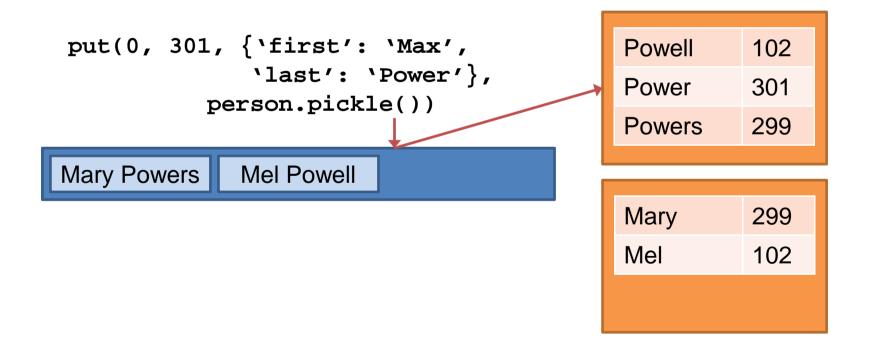
#### **Consistency: Create**

Powell	102
Powers	299

Mary	299
Mel	102

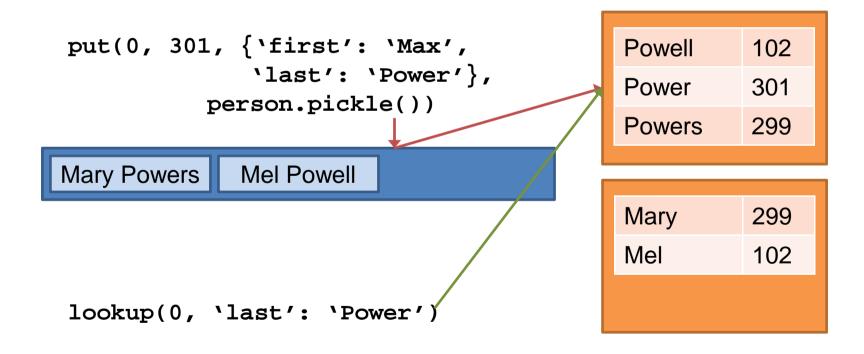
#### **Consistency: Create**

- Insert index entries before writing object
  - O What happens if a lookup happens in the meantime?



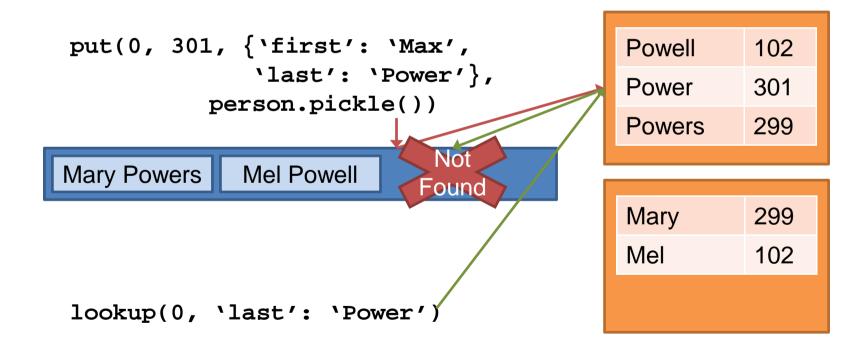
#### **Consistency: Concurrent Lookup**

Concurrent ops ignore inconsistent entries



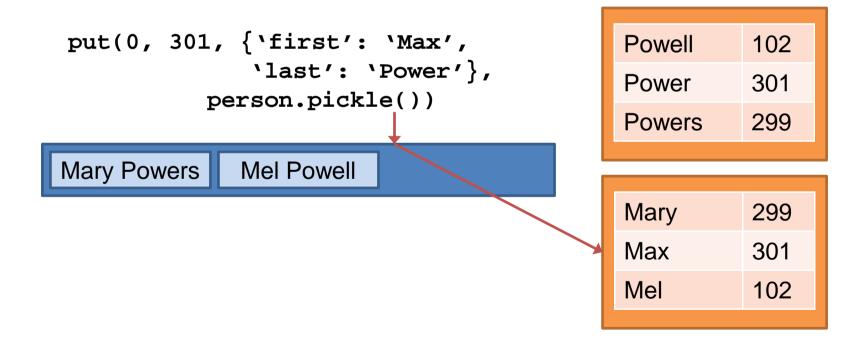
#### **Consistency: Concurrent Lookup**

Concurrent ops ignore inconsistent entries



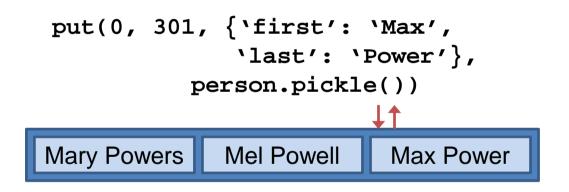
#### **Consistency: Create (continued)**

Insert index entries before writing object



#### **Consistency: Create**

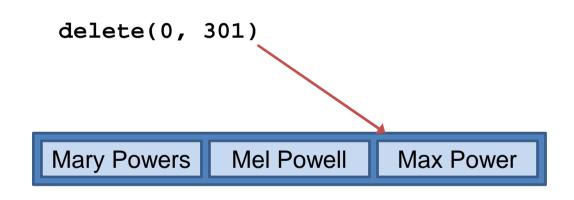
Put completes; index entries now valid



Powell	102
Power	301
Powers	299
Powers	299

299
301
102

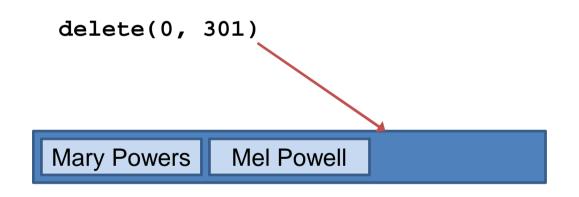
## **Consistency: Delete**



Powell	102
Power	301
Powers	299
1 011010	200

299
301
102

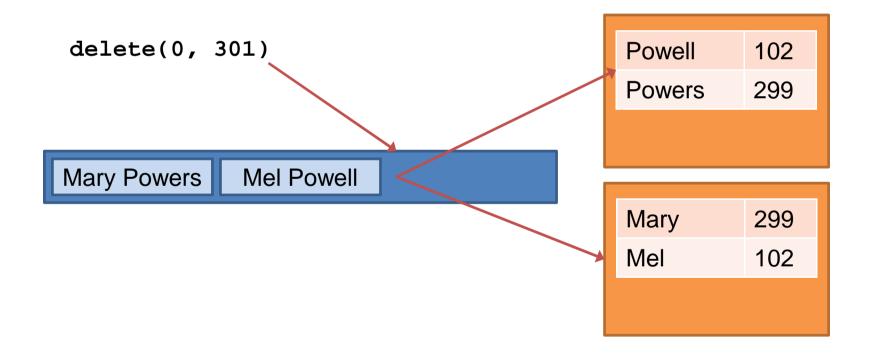
## **Consistency: Delete**



102
301
299

299
301
102

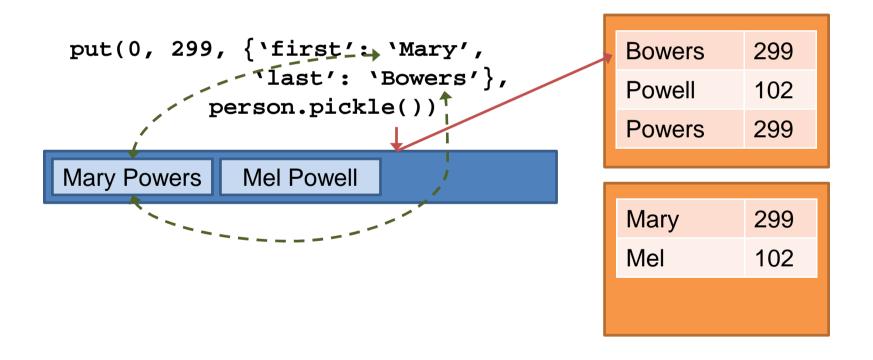
### **Consistency: Delete**



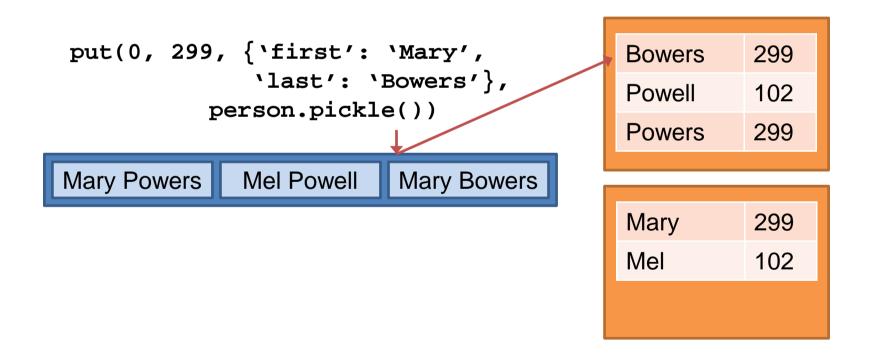
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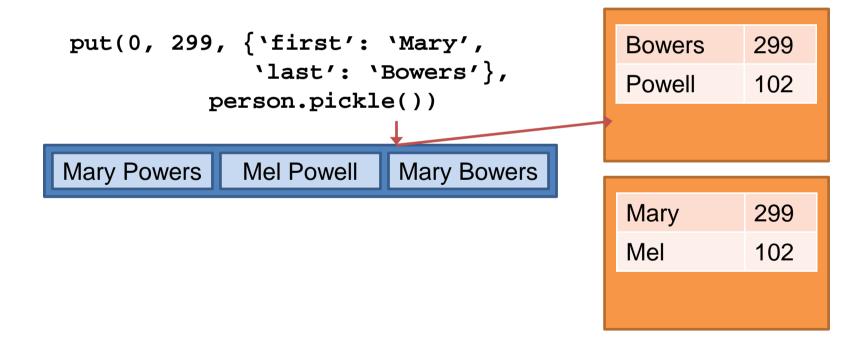
- Compare previous index entries
  - o Insert new value if updated



- Commit by writing the new value
  - Old index entries ignored by lookup since inconsistent



• Cleanup old, inconsistent entries



#### **Consistency: Thoughts**

- Low-latency gives simplified consistency
- Turn atomic puts into atomic index updates
  - All index updates for an object go through master
  - Index entries invalid until corresponding put completes

#### **Index Recovery**

- Problem: Unavailable until indexes recover
  - Many requests will be lookups
  - These will block unless indexes are recovered
- Rebuild from other masters?
  - TODO why this fails
  - o TODO Doesn't fail with sharding?
- Rebuild from backups?
  - o TODO

### **Index Recovery: Sharding**

#### Split on object ID

- Can always co-locate index with data
- Index chunk at most 320 MB
- Each new master can rebuild in a fraction of a second

#### Split on search key

- o Entire shards composed only of index data
- o At most 640 MB apiece
  - 0.6s to gather data, fraction of a second to rebuild
- Part or all of 640 MB may come from shards in recovery
  - -0.6s + 0.6s = 1.2s upper bound

#### **Index Recovery: Replication**

- Idea: Replicate indexes once in RAM
  - Threat is only to availability, not data loss
- Idea: Only preserve the shape of the index
  - The search keys are stored in the log
- TODO

## **Index Recovery: Logging**

• TODO

#### **Summary**

- Apps provide search keys explicitly on put
- Partition indexes on search key for easy lookup/enumeration
- Atomic indexes from atomic puts
- Fast index recovery for high-availability