

Low Latency RPCs

RAMCloud Design Review, April 1, 2010

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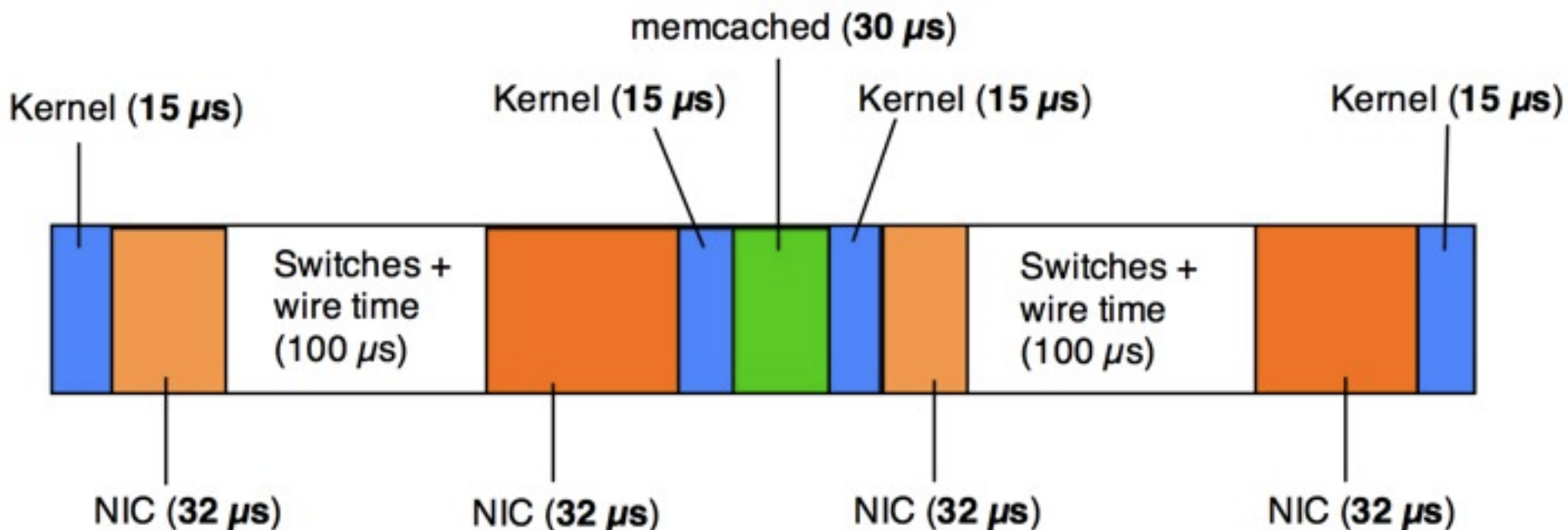


Overview

	1985	2010	Improvement
CPU Speed	12 Mhz	4 Ghz	333 x
Bandwidth	10 Mbps	10 Gbps	1000 x
Latency	2 ms	500 μ s	4 x

- Goal: **5-10 μ s** RPCs
- Experimental result: **11 μ s** RTT
- Three parts: current sources of latency, experimental results, RPC system

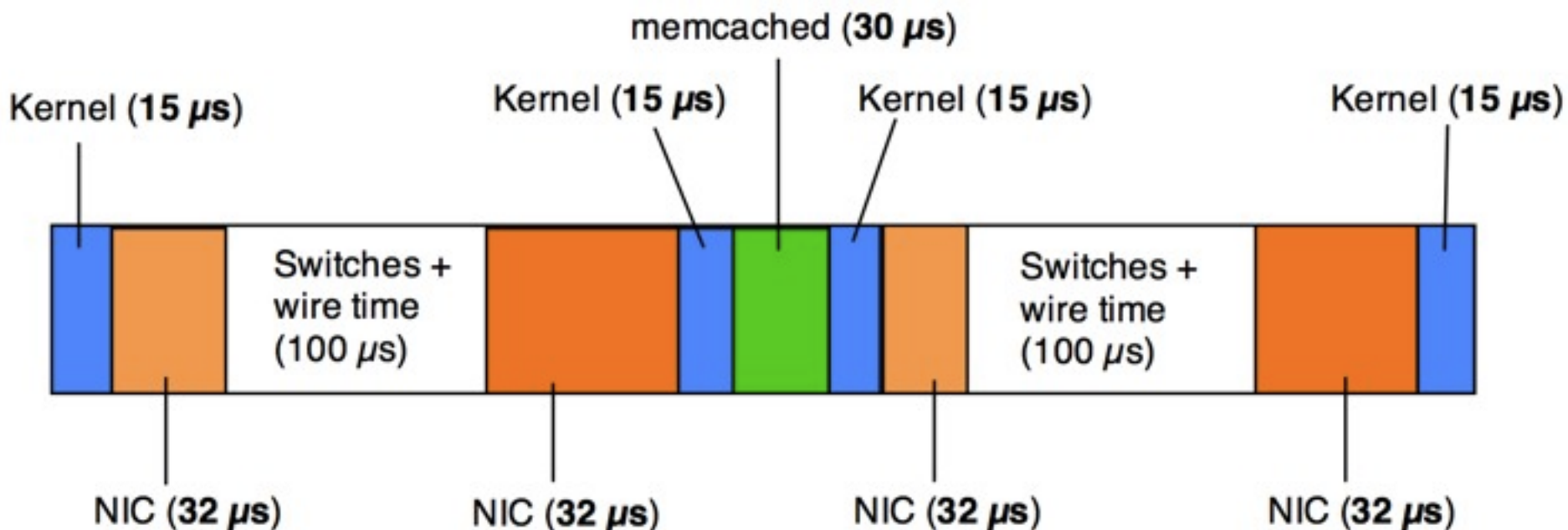
Baseline Performance



- **Experiment Setup:**

- Intel Xeon - 3.4 Ghz
- Intel 82541GI Gigabit NICs
- Standard Linux Kernel with UDP
- Switches + wire time
 - Estimated using a typical data center
 - 10 switches

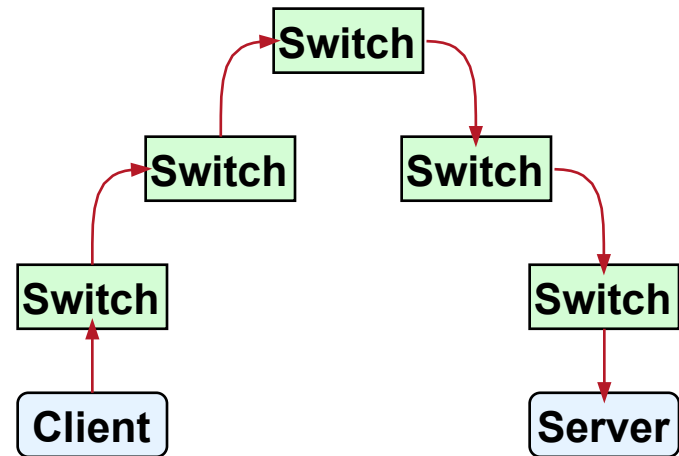
Baseline Performance



- **Total time: ~ 400 μs**
- **Main sources of latency:**
 - Switch + wire time: 200 μs
 - NIC: 128 μs
 - Kernel: 60 μs
 - memcached: 30 μs

Causes: Data center network time

- **Network latency: 150 - 300 μs**
 - 10 - 30 μs per switch, 5 switches each way
- **Latest Arista product:**
 - 0.9 μs per switch
 - Need cut-through routing, congestion management
- **Hoping for help!**
 - Not RAMCloud's goal



Causes: NIC Hardware

- **Most hardware is designed for throughput, not latency**
- **Interrupt coalescing/throttling: ~ 64 μ s one way**
 - Design the NIC to avoid live lock, and to lower CPU utilization
 - Optimize for bandwidth
 - Default setting!

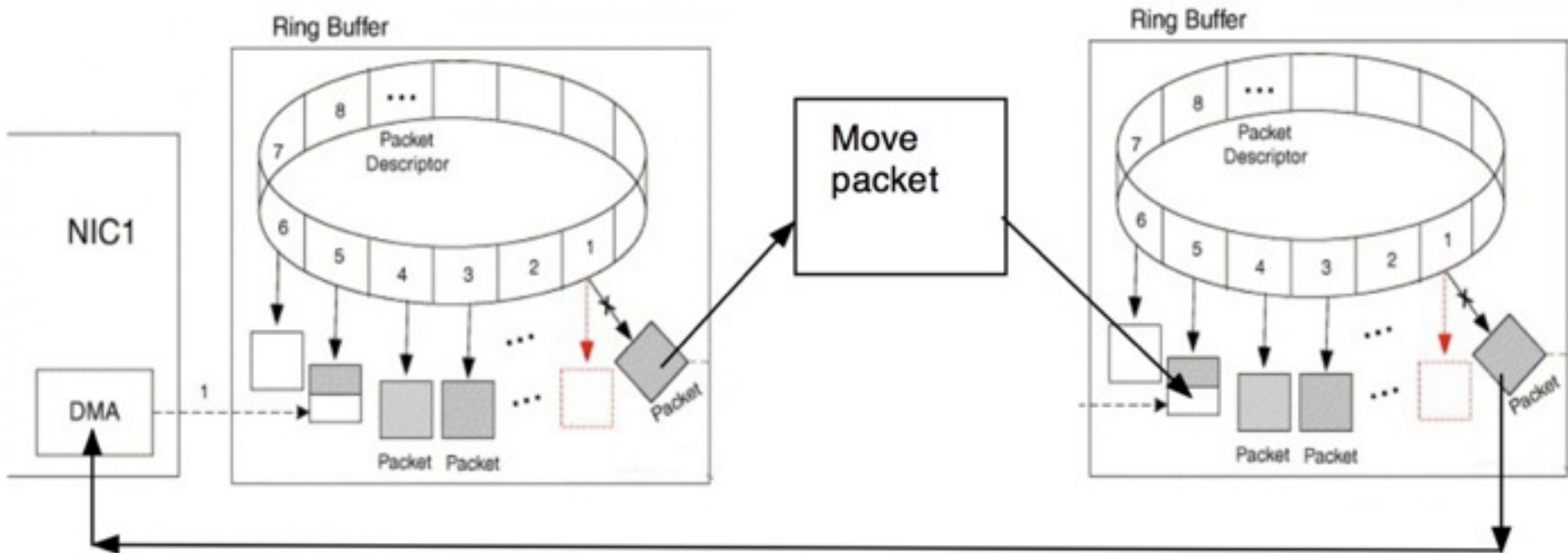
Causes: Software

- **Kernel network stack**
 - Packet takes **15 μ s** to bubble through the kernel (each way)
 - **60 μ s** of overhead per RTT!
- **Protocol overhead**
 - TCP is inherently slow
 - requires a lot of processing and state
 - IP: options may add processing time
- **Unnecessary intermediate copies**
 - From user-space to kernel
- **CPU scheduling/preemption**
- **Context switches**

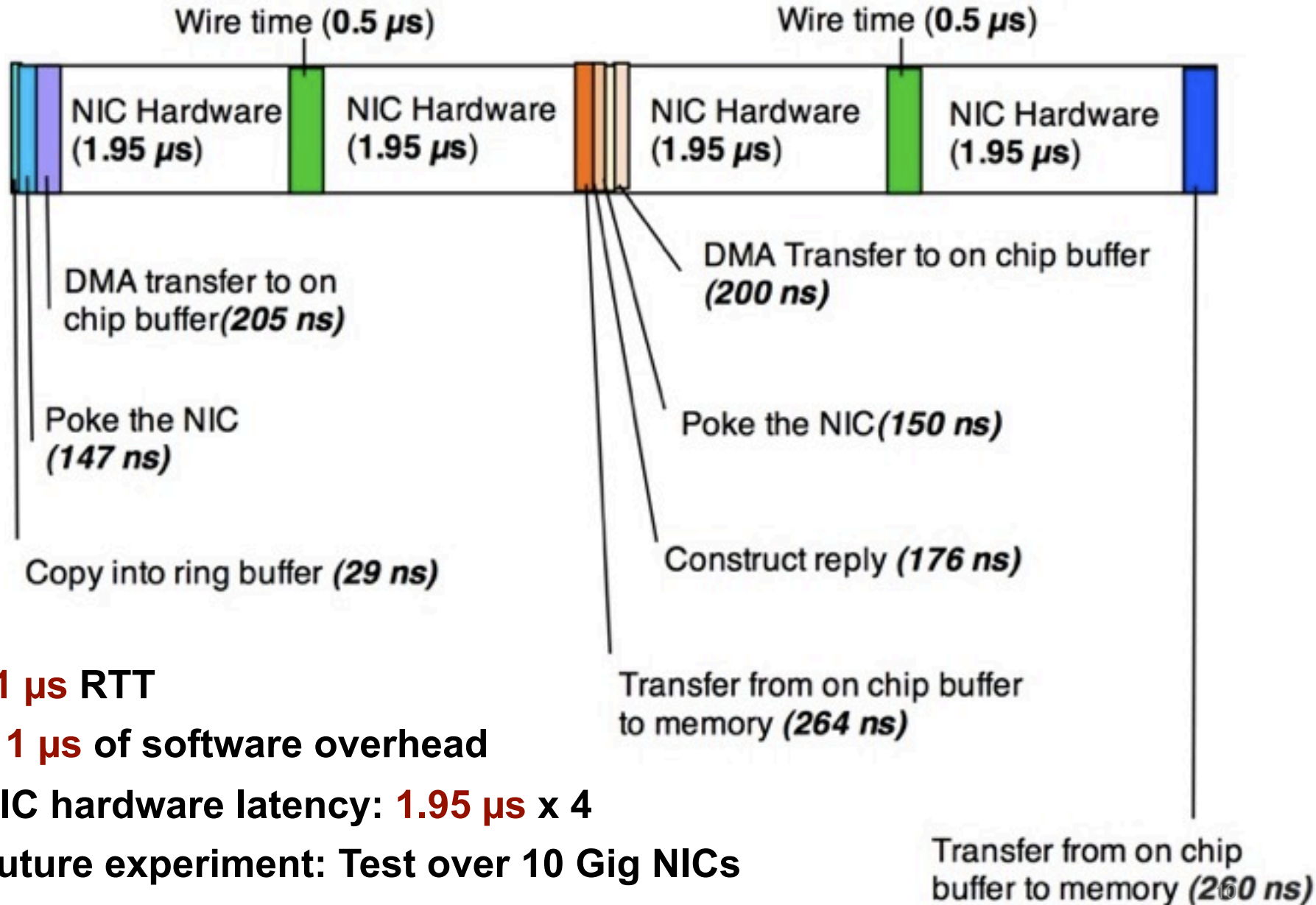
Radical Experiment

- **Part 1: Tune the NIC**
 - Turn off interrupt coalescing
 - Poll the NIC with a dedicated core, no interrupts!
- **Part 2: Rip out unnecessary layers of software**
 - Map the NIC directly into user space
 - User software can access NIC's registers and ring buffers
 - Eliminate networking layer
 - Avoids unnecessary copying
 - No kernel/context switching overhead
- **Part 3: Eliminate protocol overhead**

NIC Ring Buffers



Experimental Result



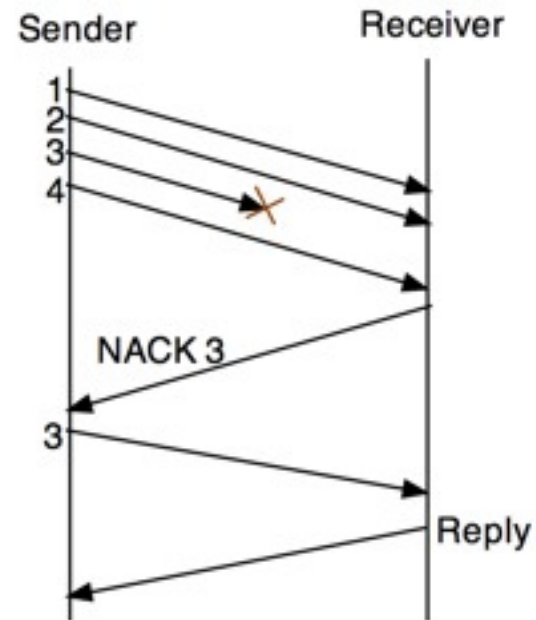
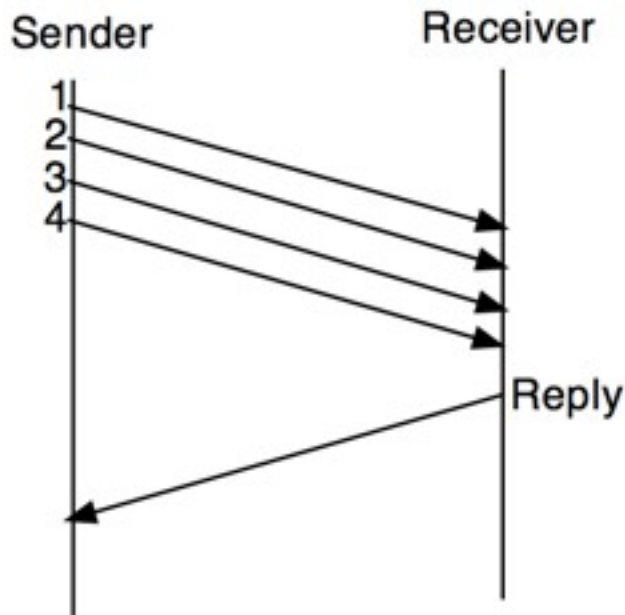
- **11 μs RTT**
- **< 1 μs of software overhead**
- **NIC hardware latency: 1.95 μs x 4**
- **Future experiment: Test over 10 Gig NICs**

RPC System

- **Build a real system**
 - As fast as weird experimental version?
- **Requirements:**
 - Reliability
 - Handles messages larger than 1 frame
 - Retain single copy
 - From ring buffer to server's log (on receive)

RPC System

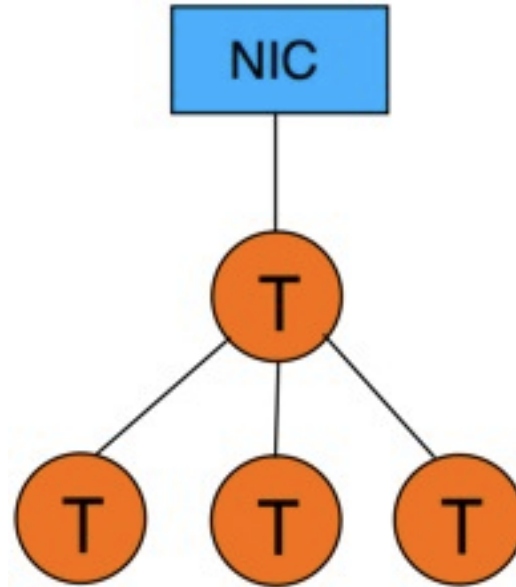
- **The reply is the ACK for most RPCs**
 - RPCs are so fast that it makes no sense to ACK fragments
- **Blast protocol**
 - Send all fragments of an RPC at once, without waiting for ACK
 - Selective NACKs
 - Too slow to retransmit the whole packet



Threading model

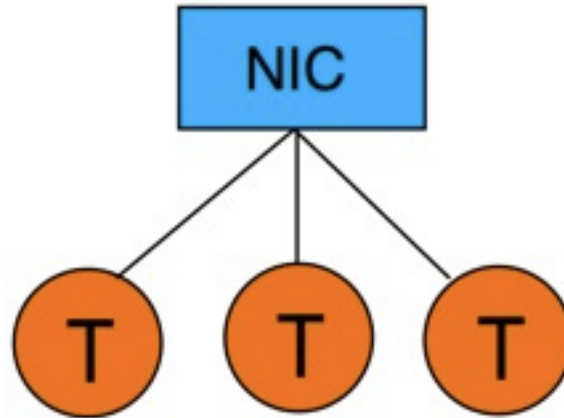
- **Increased parallelism:**
 - More cores per chip
 - More threads per core
- **Use multiple threads to increase throughput**
 - Associated dispatching/synchronization overheads
- **On server, how to distribute requests among available worker threads?**
- **Several possible designs**

Threading Model



- **Single NIC driver thread**
 - Multiplexes requests among worker threads
- **Intelligent multiplexing**
- **IPC: Shared memory regions**

Threading Model



- **Faster if we pass around the NIC?**
- **Needs locking around the NIC**

Threading Model



- **Single threaded**
- **Avoid dispatching/synchronization costs**
- **Lowest latency?**

RPC API

- **Asynchronous API:**

- Can have multiple outstanding RPCs
- Can be used by master to communicate with backups
- Can be used by client to perform multiple operations in parallel

```
rpc1.startRPC(backup1, payload);  
rpc2.startRPC(backup2, payload);  
rpc3.startRPC(backup3, payload);
```

```
// do_other_work()
```

```
Buffer *reply1 = rpc1.getReply();  
Buffer *reply2 = rpc2.getReply();  
Buffer *reply3 = rpc3.getReply();
```

- **Broadcast**

- Needed for some parts of the system: recovery, etc
- Support in RPC layer or on top of it?

Conclusion

- **Experimental fast RPCs: 11 μ s**
 - Rip out unnecessary software layers
 - NIC Hardware: 2.45 μ s x 4
- **Software overheads < 1 μ s**
 - But in an impractical ways
- **Need help with NIC and switches**
- **Early RPC system design**

Discussion

- **Is 5-10 μ s achievable? Is it worthwhile?**
- **Threading model: event based vs worker threads**
- **Should we limit the size of an RPC?**