RAMCloud 1.0

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(with Arjun Gopalan, Ashish Gupta, Ankita Kejriwal, Collin Lee, Behnam Montazeri, Diego Ongaro, Seo Jin Park, Henry Qin, Mendel Rosenblum, Stephen Rumble, and Ryan Stutsman)



Overview

• **RAMCloud project in transition**

• Phase 1 complete:

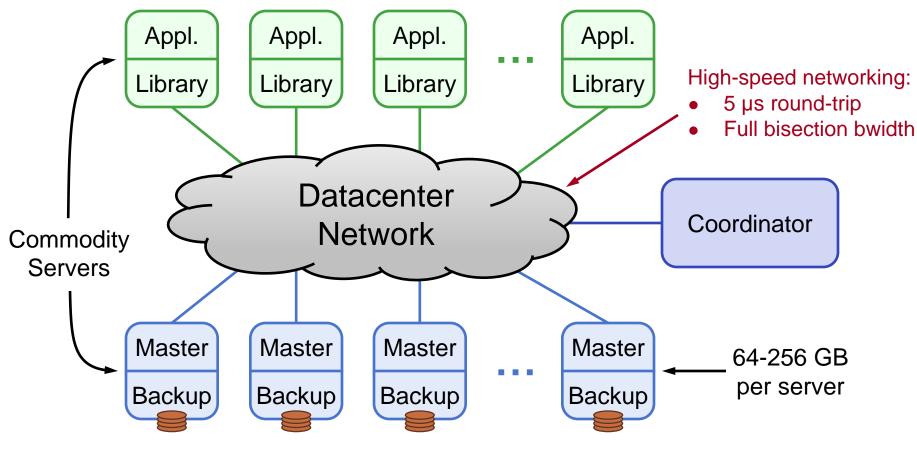
- RAMCloud 1.0 released
 - Key-value store
 - Memory management
 - Fast crash recovery
- First PhD students graduating

• Phase 2 starting up:

- New projects:
 - Higher-level data models
 - Datacenter RPC revisited
- Many new students

RAMCloud Architecture

1000 – 100,000 Application Servers



1000 – 10,000 Storage Servers

Data Model: Key-Value Store

(Only overwrite if

version matches)

• Basic operations:

- read(tableId, key)
 => blob, version
- write(tableId, key, blob)
 => version
- delete(tableId, key)

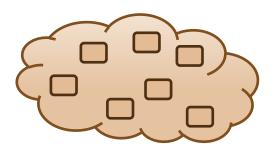
Other operations:

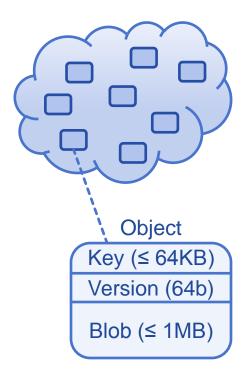
- cwrite(tableId, key, blob, version)
 => version
- Enumerate objects in table
- Efficient multi-read, multi-write
- Atomic increment

• Not currently available:

- Atomic updates of multiple objects
- Secondary indexes

Tables





Performance

Using Infiniband networking

- 24 Gb/sec effective bandwidth
- Kernel bypass
- Can also use other networking, but slower

• Reads:

- 100B objects: 5µs
- 10KB objects: 10µs
- Single-server throughput (100B objects): 700 Kops/sec.
- Small-object multi-reads: 1-2M objects/sec.

• Writes:

- 100B objects: 15µs
- 10KB objects: 40µs

Steve Rumble's Dissertation

• How to manage objects in DRAM?

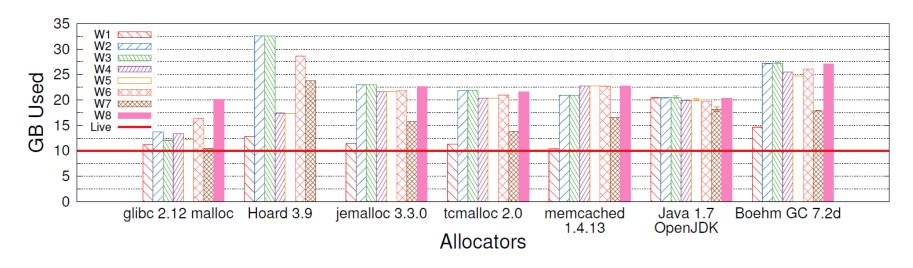
- High write performance
- High memory efficiency (80-90% utilization)
- Uniform log structure for all info (both DRAM and disk)
 - Log cleaner => incremental generational garbage collector

• Innovative aspects:

- 2-level cleaning (different policies for DRAM, disk)
- Parallel cleaning (hides cost of cleaning)
- Improved LFS segment selection formula

Paper in FAST 2014, dissertation nearing completion; Steve is working at Google Zurich

Existing Allocators Waste Memory

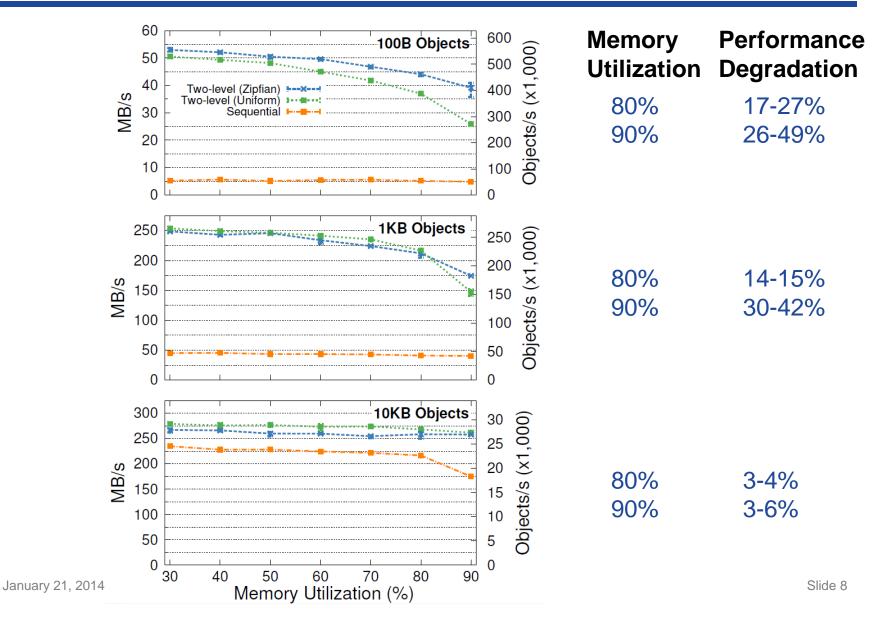


• Allocators waste memory if workloads change:

- E.g., W2 (simulates schema change):
 - Allocate 100B objects
 - Gradually overwrite with 130B objects

• All existing allocators waste at least 50% of memory under some conditions

Client Write Throughput



Ryan Stutsman's Dissertation

Durability and availability for data in DRAM

• Fault-tolerant log for each master:

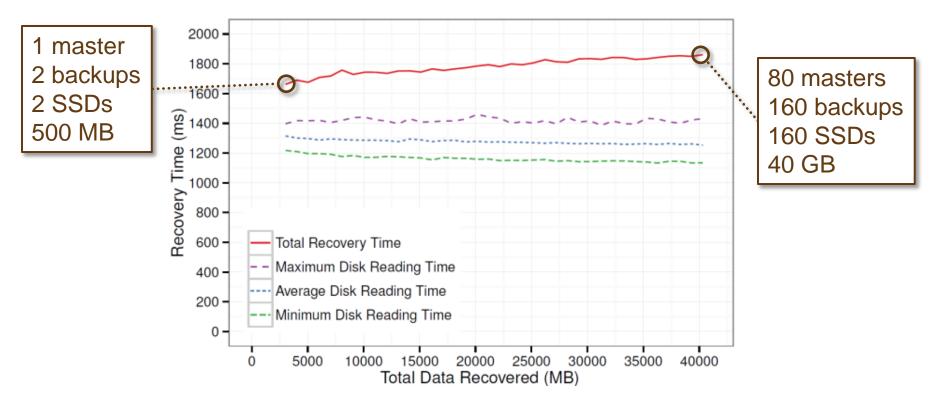
- Decentralized log management
- Finding log after crashes
- Restoring redundancy after backup crashes

• Fast crash recovery:

- Use thousands of servers concurrently to reload data from a crashed master
- 1-2 second recovery
- Handling simultaneous master/backup failures
- Dissertation filed December 2013
 Ryan is now a post-doc at Microsoft Research

January 21, 2014

Recovery Scalability



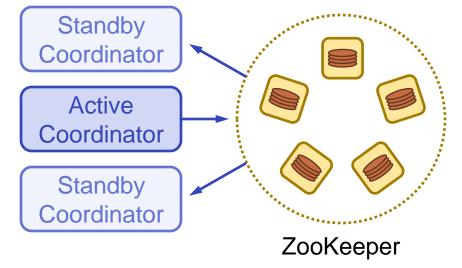
• Will improve with newer machines

- More cores (our nodes: 4 cores)
- More memory bandwidth (our nodes: 11 GB/sec)

• Bottom line: recover 500 MB/sec/server

Coordinator Crash Recovery

- Active/standby model
- Use replicated external storage for configuration data:
 - Cluster membership
 - Table metadata
- Distributed "commit" mechanism:
 - Record intent in storage
 - Notify relevant servers
 - (Eventually) mark storage "completed"
 - During restart, find and finish uncompleted ops



Missing from RAMCloud 1.0

• Table configuration management:

- Tablets not moved after creation
- RAMCloud has mechanisms for splitting, migrating tablet
- No policy code yet

• Crude reconfiguration: crash server!

Tablets split among many other servers

New Project: Data Model

- Goal: higher-level data model than just key-value store:
 - Secondary indexes
 - Transactions spanning multiple objects and servers
 - Graph-processing primitives (sets)

• Can RAMCloud support these without sacrificing

- Latency?
- Scalability?
- First project: secondary indexes (SLIK) (Arjun Gopalan, Ashish Gupta, Ankita Kejriwal)
 - Design complete, implementation underway
 - Ankita will discuss design issues

New Work: Datacenter RPC

Complete redesign of RAMCloud RPC

General purpose (not just RAMCloud)

• Latency:

- Analyze, reduce latency
- Explore alternative threading strategies
- Optimize for kernel bypass

• Scale:

- Support 1M clients/server (minimal state/connection)
- Congestion control: reservation based?

Initial projects underway (Behnam Montazeri, Henry Qin)

- Analyze latency of existing RPCs
- Support for pFabric, SolarFlare NIC

Other Projects

- Raft: new consensus protocol (Diego Ongaro)
- Graph processing on RAMCloud (Jonathan Ellithorpe)
- Using a rule-based approach for distributed, concurrent, fault-tolerant code (Collin Lee, Ryan Stutsman)

Conclusion

- We now have a usable system
- Still many open research problems
- Real usage should generate additional information, ideas