Fast Crash Recovery in RAMCloud

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SEDCL Retreat June 3, 2011

Durability and Availability

- Goals:
 - No impact on performance
 - Minimum cost, energy
- Replicate in DRAM of other masters?
 - 3x system cost, energy
 - Still have to handle power failures
 - Replicas unnecessary for performance

Durability and Availability

• RAMCloud approach:

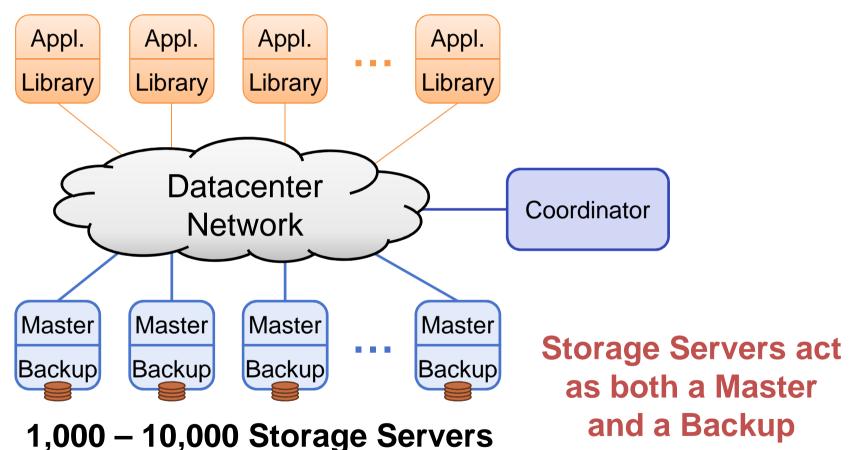
- 1 copy in DRAM
- Backup copies on disk/flash: durability ~ free!

Problems

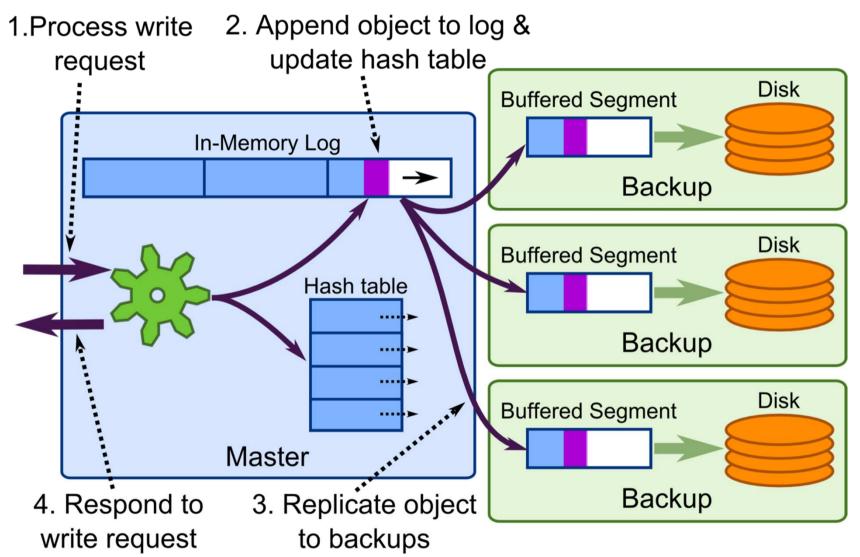
- Synchronous disk I/O's during writes?
 - Buffered logging
- Data unavailable after crashes?
 - Fast recovery

RAMCloud Architecture

1,000 – 100,000 Application Servers



Buffered Logging



Buffered Logging

• No disk I/O during write requests

- But must guarantee buffered data durability:
 - DIMMs with built-in flash backup?
 - Caches on enterprise disk controllers?
 - Per-server battery backups?

• 8 MB I/Os minimize disk latency overhead

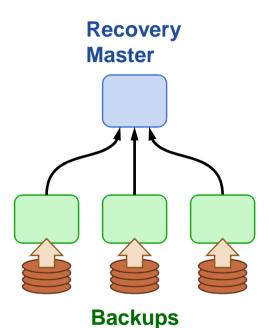
- Achieves 90% disk bandwidth
- Helpful for reading during fast recovery as well
- Master's memory also log-structured
 - Log cleaning ~ generational garbage collection

Master Crash Recovery

- Problem: Data unavailable after a master crash
- Goal: 1 2 second recovery
 - Applications just see "hiccups"
 - Good enough for "continuous availability?"
- Solution: Harness system scale

Recovery, First Try

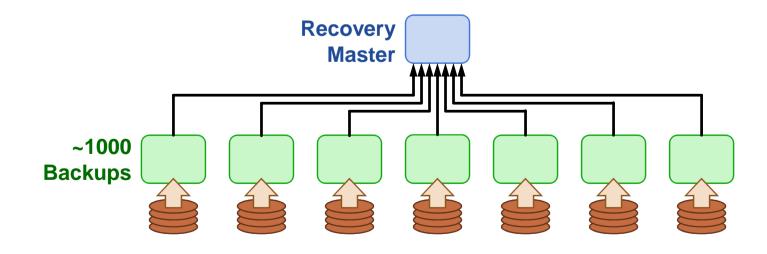
- Master chooses backups statically
 - Each backup stores entire log for master
- Crash recovery:
 - Choose recovery master
 - Backups read log info from disk
 - Transfer logs to recovery master
 - Recovery master replays log
- First bottleneck: disk bandwidth:
 - 64 GB / 3 backups / 100 MB/sec/disk
 ≈ 210 seconds
- Solution: more disks (more backups)



Recovery, Second Try

• Scatter logs:

- Each log divided into 8MB segments
- Master chooses different backups for each segment (randomly)
- · Segments scattered across all servers in the cluster
- Crash recovery:
 - All backups read from disk in parallel
 - Transmit data over network to recovery master

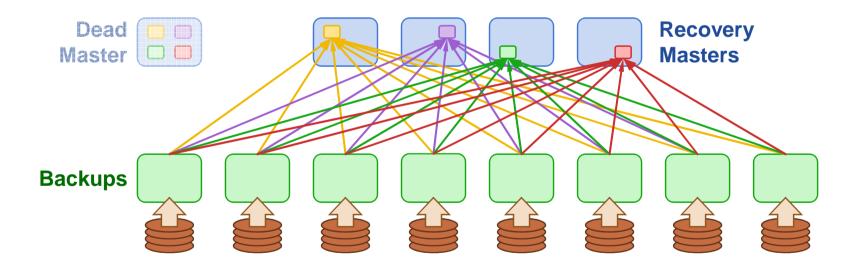


Scattered Logs

- Disk no longer a bottleneck:
 - 64 GB / 8 MB/segment / 1000 backups ≈ 8 segments/backup
 - 100 ms/segment to read from disk
 - 0.8 seconds to read all segments in parallel
- Second bottleneck: NIC on recovery master
 - 64 GB / 10 Gbits/second ≈ 60 seconds
 - Recovery master CPU is also a bottleneck
- Solution: more recovery masters
 - Spread work over 100 recovery masters
 - 64 GB / 10 Gbits/second / 100 masters ≈ 0.6 seconds

Recovery, Third Try

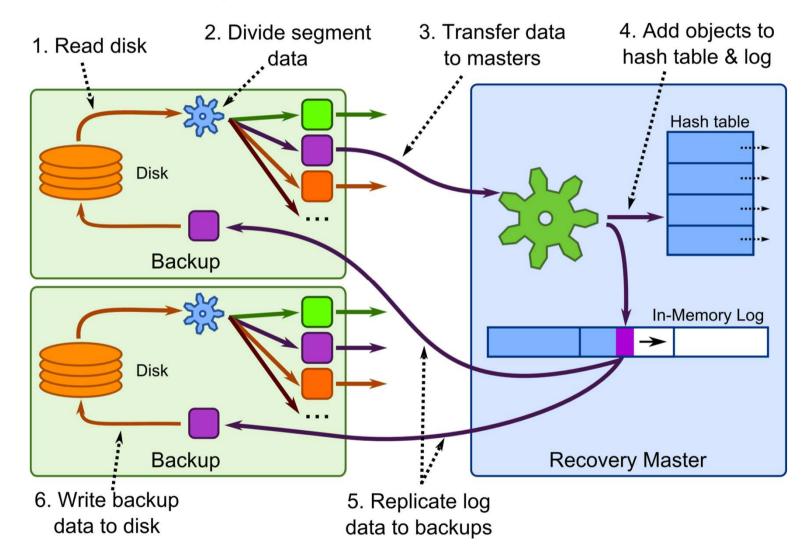
- Divide each master's data into partitions
 - Recover each partition on a separate recovery master
 - Partitions based on tables & key ranges, not log segment
 - Each backup divides its log data among recovery masters



Start of Recovery

- 1. Coordinator C gets "M is down" report
- 2. C verifies M is down
- 3. C broadcasts to cluster: "which segments do you have for M?" "begin reading and partitioning them"
- 4. C verifies no segments are missing
- 5. C notifies selected recovery masters: "recover this partition of M from these locations"

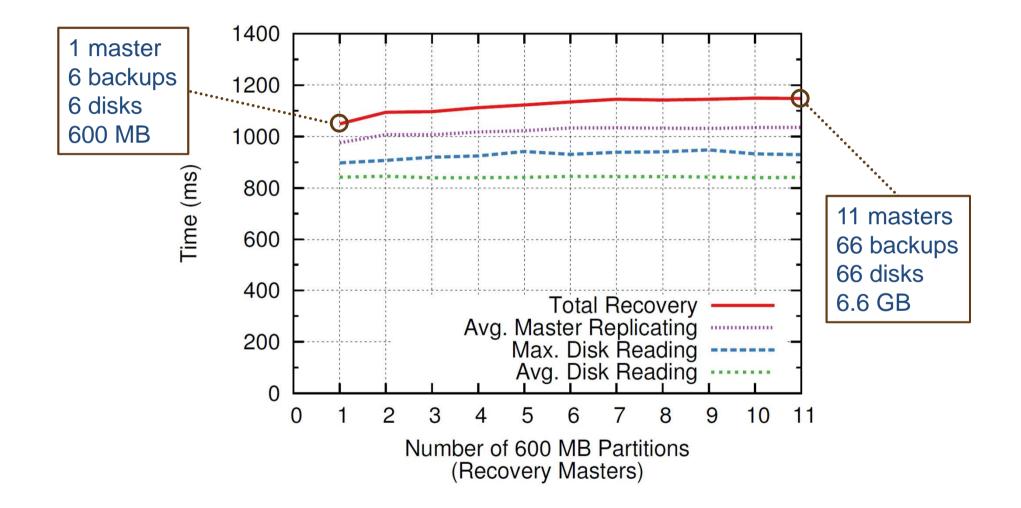
Recovery



Parallelism During Recovery

- Tightly overlapped and pipelined
 - Each recovery master can recover 400-800 MB of log per second
 - All steps are all being performed by all hosts simultaneously
- Segments can be replayed in any order
 - Possible due to version number on objects
 - Recovery masters can work independently and distribute load over backups
 - Recovery masters can replay any data when it becomes available

Recovery Scalability



TODO

- Restoring locality
- Recovering Backups
- Cold boot

Recovering Backups

- In parallel with master recovery
- Each master recreates any segments it has stored on that backup elsewhere
- Expect to re-replicate 8 segments
 - 1,000 machines, 64 GB per master
 - Just need to buffer on a new backup
 - Fast compared to master recovery

Recovery Summary

Before Crash

- Scatter log data
- Balance partitions
 - Steve's talk

On Crash

- Detect failure
- Find log data
- Check log integrity
- Select new masters
- Replay log data
- Recreate backups